FastLAS: Scalable Inductive Logic Programming, incorporating Domain-specific Optimisation Criteria

Mark Law, Alessandra Russo, Elisa Bertino, Krysia Broda and Jorge Lobo

February 11, 2020

Inductive Logic Programming

Input: background knowledge B and examples E^+ and E^- .

Output: a hypothesis H s.t.

```
1. \forall e^+ \in E^+ : B \cup H \models e^+
```

$$2. \ \forall e^- \in E^- : B \cup H \not\models e^-.$$

Inductive Logic Programming

Input: background knowledge B and examples E^+ and E^- .

Output: a hypothesis H s.t.

1. $\forall e^+ \in E^+ : B \cup H \models e^+$ 2. $\forall e^- \in E^- : B \cup H \not\models e^-$.

The paper contains two main contributions:

- FastLAS, an ASP-based ILP system which is scalable w.r.t. the size of the hypothesis space.
- A new approach to supporting domain specific optimisation criteria in the search.

Optimisation Criteria

In some domains, false negatives are more dangerous than false positives.



Event detection



Medical diagnosis

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Depending on the domain, different optimisation criteria may be required!

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Many other optimisation criteria exist:

- ▶ The number of rules in *H*.
- ▶ The number of ground instances of rules in *H*.
- ► The number of variables in H.
- ▶ The number of answer sets of $B \cup H$.

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A general scoring function takes as input a learning task T and a hypothesis H and returns a score in $\mathbb{R}_{>0}$.

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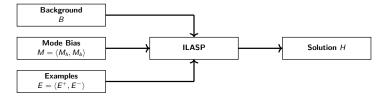
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FastLAS currently supports *decomposable* scoring functions. These can be evaluated on each rule independently.

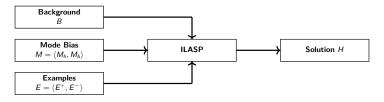
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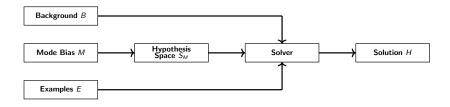


FastLAS has two major restrictions:

- ► All programs have a single answer set.
- ▶ No predicate in M_h occurs in M_b or any rule in T.

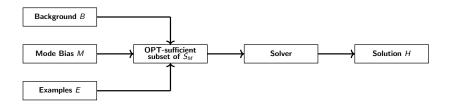
FastLAS: Overview

ILASP begins by computing the hypothesis space S_M , which contains every rule that is compatible with the mode bias M.



FastLAS: Overview

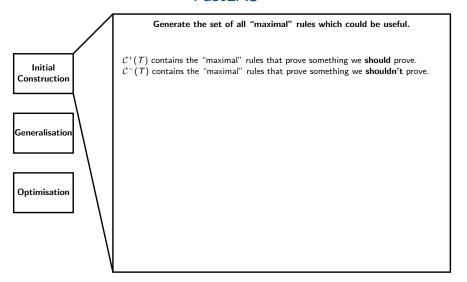
FastLAS instead begins by computing an *OPT-sufficient* subset of the hypothesis space, which is often significantly smaller than S_M .

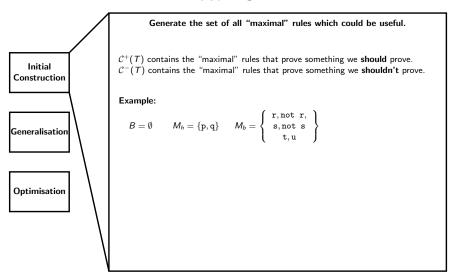


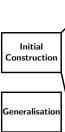
Initial Construction

Generalisation

Optimisation







Optimisation

Generate the set of all "maximal" rules which could be useful.

 $\mathcal{C}^+(T)$ contains the "maximal" rules that prove something we **should** prove. $\mathcal{C}^-(T)$ contains the "maximal" rules that prove something we **shouldn't** prove.

Example:

$$B = \emptyset$$
 $M_b = \{p, q\}$ $M_b = \left\{ egin{array}{ll} r, not & r, \\ s, not & s \\ t, u \end{array} \right\}$

Find
$$H$$
 s.t.
$$\left\{ \begin{array}{l} H \models q, & H \not\models p, \\ H \cup \{r. \ t.\} \models p, & H \cup \{r. \ t.\} \not\models q, \\ H \cup \{r. \ u.\} \models p, & H \cup \{r. \ u.\} \not\models q \end{array} \right.$$



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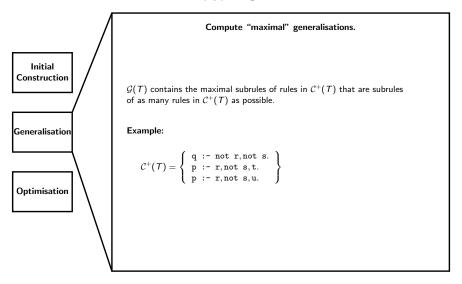
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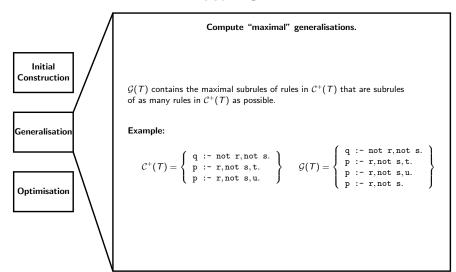
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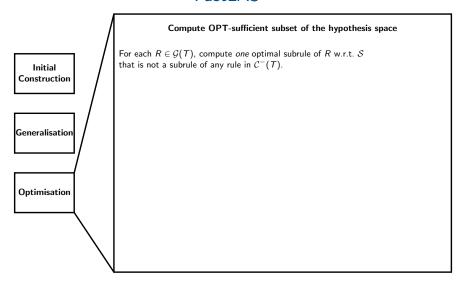
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$$\mathcal{C}^+(\mathcal{T}) = \left\{ \begin{array}{l} q \text{ :- not r,not s.} \\ p \text{ :- r,not s,t.} \\ p \text{ :- r,not s,u.} \end{array} \right\} \quad \mathcal{C}^-(\mathcal{T}) = \left\{ \begin{array}{l} p \text{ :- not r,not s.} \\ q \text{ :- r,not s,t.} \\ q \text{ :- r,not s,u.} \end{array} \right\}$$









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Compute OPT-sufficient subset of the hypothesis space

For each $R\in\mathcal{G}(T)$, compute *one* optimal subrule of R w.r.t. \mathcal{S} that is not a subrule of any rule in $\mathcal{C}^-(T)$.

Example:

$$\mathcal{G}(T) = \left\{ \begin{array}{l} q : - \text{ not } r, \text{not } s. \\ p : - r, \text{not } s, t. \\ p : - r, \text{not } s, u. \\ p : - r, \text{not } s. \end{array} \right\} \quad \mathcal{C}^{-}(T) = \left\{ \begin{array}{l} p : - \text{ not } r, \text{not } s. \\ q : - r, \text{not } s, t. \\ q : - r, \text{not } s, u. \end{array} \right\}$$

Optimise with S_{len} :



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 $\mathcal{O}(T,\mathcal{S})$ is *OPT-sufficient*. Hence, FastLAS is guaranteed to return an optimal solution w.r.t. any decomposable scoring function!

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The full hypothesis space (computed by ILASP) contains 72 rules.

Evaluation

Sentence Chunking

In (Kazmi et al. 2017), the Inspire system was evaluated on a sentence chunking dataset (Agirre et al. 2016), which contains examples of how sentences should be *chunked*.

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| System | F_1 | Running Time | |
|---------|-------|----------------------|--|
| INSPIRE | 0.712 | >1800s in some cases | |
| ILASP3 | 0.777 | 1051.4s | |
| FastLAS | 0.768 | 4.5s | |

We evaluated FastLAS on access control datasets.

▶ We used three scoring functions, S_{len} , S_{cov} and S_{uni} , which encouraged learning progressively more general hypotheses.

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Learning rules for accept:

| Recall | Precision |
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| 0.905 | 0.951 |
| 0.892 | 0.949 |
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Learning rules for reject:

| Scoring function | Recall | Precision |
|------------------------------|--------|-----------|
| $\mathcal{S}_{\mathit{len}}$ | 0.974 | 0.935 |
| \mathcal{S}_{cov} | 0.969 | 0.941 |
| \mathcal{S}_{uni} | 0.966 | 0.965 |

Conclusion

FastLAS is a new ASP-based ILP system.

- Far more scalable than ILASP.
- Supports domain-specific optimisation criteria, through scoring functions.

Future research directions include:

- Lifting current restrictions, to allow recursion, non-observational predicate learning, predicate invention and programs with multiple answer sets.
- Non-decomposable scoring functions.

Backup Slides

Sentence Chunking Results

100 examples:

| System | F_1 | Running Time |
|---------|-------|--------------|
| INSPIRE | 0.733 | _ |
| ILASP3 | 0.757 | 210.4s |
| FastLAS | 0.751 | 0.909s |

500 examples:

| System | F_1 | Running Time |
|---------|-------|--------------|
| INSPIRE | 0.712 | _ |
| ILASP3 | 0.777 | 1051.4s |
| FastLAS | 0.768 | 4.5s |

CAVIAR

The CAVIAR dataset (generated by the EC Funded CAVIAR project/IST 2001 37540) consists of manually annotated video streams.





The task is to detect events by learning initiating/terminating conditions. We compared FastLAS to OLED (Katzouris et al. 2016) and ILASP3 (Law et al. 2015) on the task of detecting people meeting.

CAVIAR Results

| System | F_1 | Running Time |
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| OLED | 0.792 | 107s |
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For CAVIAR, ILASP had a hypothesis space with 3370 rules. FastLAS had over 2⁴⁴ rules.

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Due to the larger search space, FastLAS achieved an $\it F_{\rm 1}$ score of 0.923 compared to ILASPs 0.842.

Policy Learning II

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| Method | Resource 25993 | Resource 4675 | Resource 75078 | Resource 79092 |
|---------------------|----------------|---------------|----------------|----------------|
| Rhapsody | 0.04 | 0.10 | 0.10 | 0.04 |
| CTA | 0.04 | 0.12 | 0.10 | 0.04 |
| FastLAS: S_{len} | 0.02 | 0.04 | 0.02 | 0.02 |
| FastLAS: S_{cov} | 0.02 | 0.04 | 0.02 | 0.02 |
| FastLAS: S_{uni} | 0.01 | 0.04 | 0.02 | 0.02 |
| FastLAS: S_{UF_1} | 0.07 | 0.10 | 0.11 | 0.05 |



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